

# Executing code on columnar data

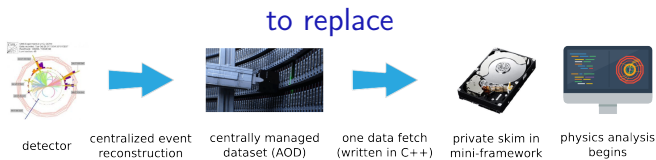
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I'm working on a query language and database server to aggregate large samples of HEP data on the fly.

**Purpose:** to eliminate the need for private skims in most situations.



Collaborating with Jin Chang and Igor Mandrichenko on the server.

The query language, Femtocode, plays a similar role as TTreeFormula:

- ▶ a high-level language for the physicist
- ▶ usually for filling a histogram (so query responses are small)
- ▶ but generally useful for transforming one dataset into another.

However, it's a full-fledged language with assignments and user-defined functions, so that it can encompass a larger part of the data analysis.

(I've examined SQL, LINQ, and others, and they are not sufficient. I would use a standard if I could. Femtocode BNF has  $> 50\%$  overlap with Python BNF.)

The essential feature of FemtoCode is that it can compile complex structure-manipulations, which would ordinarily have to be performed in object-oriented code, into a series of vectorized kernels.

It operates on columns.

## Example:

```
hist = dataset.bin(100, 0, 50, ""  
    muons.map(m => sqrt(m.px**2 + m.py**2)).max()  
    """)
```

## compiles to

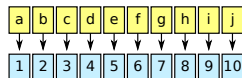
1. Compute  $\sqrt{p_x^2 + p_y^2}$  for all muons, ignoring event boundaries.
2. Find the maximum such value for each event.
3. Bin those events and fill the histogram.

## rather than

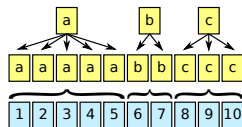
1. Loop over events:
  - 1.1 Loop over muons:
    - 1.1.1 Compute  $\sqrt{p_x^2 + p_y^2}$  for each.
  - 1.2 Fill a histogram with the maximum.

## Three types of data transformations:

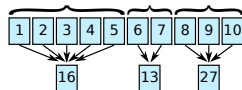
**Flat:** apply  $N$ -argument function to each element of  $N$  aligned arrays, ignoring boundaries.



**Explode:** emulate (nested) for-loops by replicating data in one array so that it becomes aligned with another array.



**Reduce:** emulate counters, sum, mean, max, etc. by combining elements of an array so that it becomes aligned with an outer level of structure.



The majority of steps in a typical calculation are flat:

```
double in[ZILLION];  
double out[ZILLION];  
  
for (int i = 0; i < ZILLION; i++)  
    out[i] = flat_operation(in[i]);
```

- ▶ Compilation with `-O3` vectorizes if possible (depends on `flat_operation`).
- ▶ Easiest form for CPU to prefetch memory and/or pipeline operations.
- ▶ Also ideal for GPU calculations.
- ▶ There is a standard for functions of this form: Numpy's `ufunc` is widely used among scientific libraries.
  - ▶ Easy way for a user to add functions to the language!

```
import ctypes, numpy, numba

libMathCore = ctypes.cdll.LoadLibrary("libMathCore.so")
chi2_ctypes = libMathCore._ZN5TMathl7ChisquareQuantileEdd # c++filt!
chi2_ctypes.argtypes = (ctypes.c_double, ctypes.c_double)
chi2_ctypes.restype = ctypes.c_double

# compile to pure-C ufunc
@numba.vectorize(["f8(f8, f8)"], nopython=True)
def chi2_ufunc(p, ndf):
    return chi2_ctypes(p, ndf)

p = numpy.random.uniform(0, 1, int(1e6)) # million random numbers
result = chi2_ufunc(p, 100) # call ufunc on all of them
# 3.22 seconds

import ROOT
result = [ROOT.TMath.ChisquareQuantile(pi, 100) for pi in p]
# 9.32 seconds
```

(Performance comparison is just to show that the ufunc computes ChisquareQuantile in C, not in Python. Simpler functions show a more dramatic difference.)



Depends critically on the way we represent structure. For the “recursive counter” method I described in the last talk,

Given:                    [ [ a b c ] [ d e f g ] ] [ [ h ] [ i j ] ]  
 Data array:                a b c                d e f g                h                i j  
 Recursive counter: 2 3                    4                    2 1                2

Calculating arbitrary explosions is solved in two cases:

- ▶ explode scalar to fit a list’s counter (35 lines of C)
- ▶ explode list to fit another list’s counter (470 lines, recursive).

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Illustration of scalar-to-list:

`xs → [1, 2, 3, 4], [], [5, 6, 7]` and `y → 100, 200, 300`

Computing

```
xs.map(x => x + y)
```

yields

```
[101, 102, 103, 104], [], [305, 306, 307]
```

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Illustration of list-to-deeper-list:

`xss` → `[[100, 200], [300, 400], [500, 600]]` and `ys` → `[1, 2, 3, 4]`

Computing

```
xss.map(xs => xs.map(x => ys.map(y => x + y)))
```

yields

```
[[ [101, 102, 103, 104], [201, 202, 203, 204]],
  [ [301, 302, 303, 304], [401, 402, 403, 404]],
  [ [501, 502, 503, 504], [601, 602, 603, 604]]]
```

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 Data array: `a b c    d e f g    h    i j`  
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Calculating arbitrary explosions is solved in two cases:

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Another illustration of list-to-deeper-list:

`xss` → `[[100, 200], [300, 400], [500, 600]]` and `ys` → `[1, 2, 3, 4]`

Computing

```
xss.map(xs => ys.map(y => xs.map(x => x + y)))
```

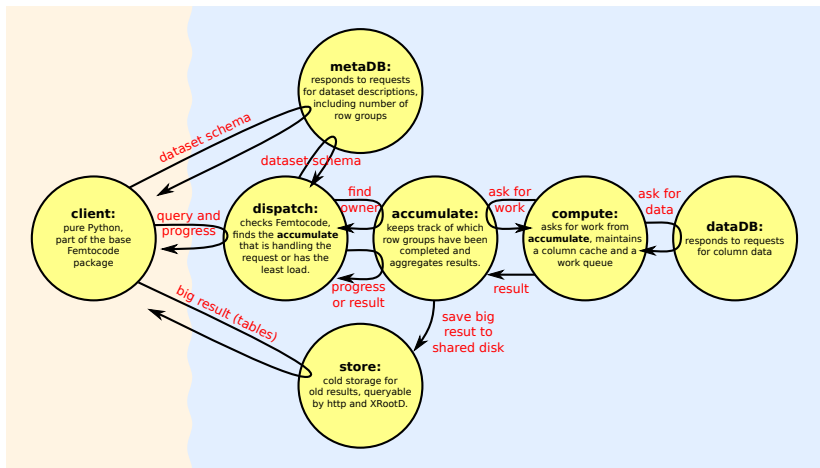
yields

```
[[[101, 201], [102, 202], [103, 203], [104, 204]],
 [ [301, 401], [302, 402], [303, 403], [304, 404]],
 [ [501, 601], [502, 602], [503, 603], [504, 604]]]
```

Haven't been implemented, but they're pretty straightforward.

Before finishing the language, we want to understand how it will fit into the server.

### Preliminary design:



If a centralized query server is going to replace private skims, it has to respond to aggregations over whole datasets in seconds.

**Purpose of early studies:** determine what performance is *possible*.

File-reading rates in events/ms per process (kHz per process), with the goal of extracting only  $p_T$ .

| particle | #/event | # branches | CMSSW | TTree::Draw() | fast reader |
|----------|---------|------------|-------|---------------|-------------|
| photon   | 2.9     | 205        | 1.14  | 435           | 769         |
| electron | 2.5     | 231        | 1.02  | 417           | 833         |
| muon     | 2.7     | 192        | 1.02  | 16.5          | 770         |
| tau      | 6.3     | 88         | 1.55  | 244           | 417         |
| jet      | 16.7    | 95         | 1.15  | 123           | 182         |
| AK8 jet  | 1.8     | 95         | 2.10  | 556           | 1000        |

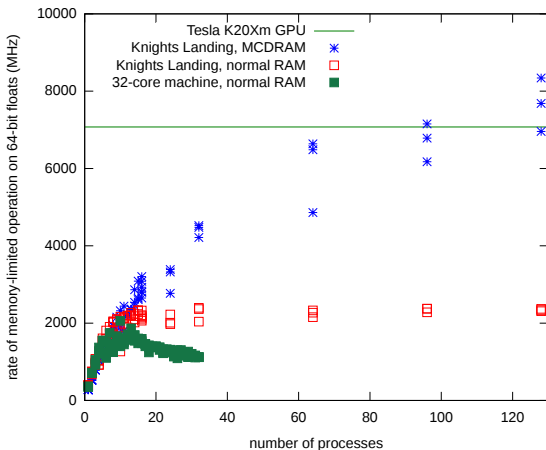
- ▶ CMSSW loads all branches to reconstruct particles as a C++ objects. Loading all branches just to cut on  $p_T$  is wasteful.
- ▶ TTree::Draw() is more streamlined, only loads required branches. (Low rate for muons is not understood.)
- ▶ “fast reader” is based on a code snippet Philippe prepared for me, using some of the same techniques as TTree::Draw().



We also plan to maintain an in-memory cache of recently used *columns*, on the supposition that the column-popularity distribution is steep enough to cause frequent cache-hits among users.

Rate for simple, flat functions on cached columns is limited only by memory bandwidth.

Could reach a peak of 7 GHz on KNL or GPU.



- ▶ I'm developing Femtocode to translate object semantics into vectorized operations as part of a project to create a fast query server.
- ▶ The “recursive counter” representation of nested structure can be exploded and reduced.
  - ▶ This representation is identical to ROOT's for depth-1 lists.
  - ▶ Any interest in extending to arbitrary split depth?
- ▶ Flat functions are
  - ▶ quick to compute,
  - ▶ extensible using Numpy's “ufunc” standard.
- ▶ For a cached query server,
  - ▶  $\sim 1$  MHz column entries is attainable for cache-misses,
  - ▶  $\sim 1$  GHz column entries is attainable for cache-hits.