#### Introduction to Quantum Computing: Qubits, Gates, and Algorithms

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# Quantum Information Science and Technology



Quantum Information Science utilizes a quantum mechanical description of nature to compute, sense, and communicate information in ways unobtainable by means based on a classical description of nature

# **Computing Development Timeline**

#### **Classical Computing (Electronic)**



Quantum computing is transitioning from scientific curiosity to technical reality.

Advancing from discovery to useful machines takes time & engineering

You must be in the game to play



18 cores

32 cores

# Quantum Worldwide (not exhaustive)



\* European Commission

#### Nascent Commercial Quantum Processors



IBM

To realize the promise of QC, we must engineer quantum systems that are robust, reproducible, and extensible.





#### Outline

#### □ Introduction

- □ Classical and Quantum Bits
- Quantum Gates and Algorithms
- Engineering Quantum Systems

#### **Classical Computer**

Fundamental logic element	"Bit" : classical bit (transistor, spin in magnetic memory, …)		
State	0 "Or" 1		
Measurement	<ul> <li><i>Discrete</i> states</li> <li>Deterministic measurement: Ex: Set as 1, measure as 1</li> </ul>		

	<b>Classical Computer</b>	Quantum Computer	
Fundamental logic element	"Bit" : classical bit (transistor, spin in magnetic memory,)	"Qubit" : quantum bit (any coherent two-level system)	
State	0 "Or" 1	$ 0\rangle \qquad Superposition: \\ \alpha 0\rangle + \beta 1\rangle \\  0\rangle \qquad ``And''  1\rangle \\  \psi\rangle = \alpha \begin{bmatrix} 1 \\ 0 \end{bmatrix} + \beta \begin{bmatrix} 0 \\ 1 \end{bmatrix}$	
Measurement	<ul> <li><i>Discrete</i> states</li> <li>Deterministic measurement: Ex: Set as 1, measure as 1</li> </ul>	<ul> <li>Superposition states</li> <li>Probabilistic measurement: Ex: If  α  =  β , 50%  0&gt;, 50%  1&gt;</li> </ul>	

Quantum computers rely on encoding information in a fundamentally different way than classical computers

#### **Classical Computer**

Fundamental logic element	"Bit" : classical bit (transistor, spin in magnetic memory, …)
Computing	<ul> <li>N bits: One N-bit state</li> <li>000, 001,, 111 (N = 3)</li> <li>Change a bit: new calculation (classical parallelism)</li> <li>000 →  ( ) →  ( 000)</li> <li>001 →  ( ) →  ( 001)</li> </ul>
	U 🔿

	Classical Computer	Quantum Computer
Fundamental logic element	"Bit" : classical bit (transistor, spin in magnetic memory, …)	"Qubit" : quantum bit (any coherent two-level system)
Computing	<ul> <li>N bits: One N-bit state</li> <li>000, 001,, 111 (N = 3)</li> <li>Change a bit: new calculation (classical parallelism)</li> </ul>	<ul> <li>N qubits: 2<sup>N</sup> components to one state <ul> <li>α 000⟩ + β 001⟩ + ··· + γ 111⟩ (N = 3)</li> </ul> </li> <li>Quantum parallelism &amp; interference</li> </ul>
	$000 \longrightarrow \boxed{2} \longrightarrow f(000)$ $001 \longrightarrow \boxed{2} \longrightarrow f(001)$	$ \begin{array}{c c} \alpha & 0 & 0 \\ \hline \end{array} & \bullet \\ \beta & 0 & 0 \\ \hline \end{array} & \bullet \\ \end{array} \rightarrow \begin{array}{c c} \alpha' & f(0 & 0 & 0 \\ \hline \end{array} & \bullet \\ \beta' & f(0 & 0 & 1 \\ \hline \end{array} & \bullet \\ \end{array} \rightarrow \begin{array}{c c} \alpha' & f(0 & 0 & 0 \\ \hline \end{array} & \bullet \\ \beta' & f(0 & 0 & 1 \\ \hline \end{array} & \bullet \\ \end{array} $

Quantum computers rely on encoding information in a fundamentally different way than classical computers

## **Classical and Quantum Bits**



 $2^{N} \rightarrow 2^{3} = 8 \rightarrow \{c_{1}, c_{2}, c_{3}, c_{4}, c_{5}, c_{6}, c_{7}, c_{8}\}$ 



Coefficients are shuttled between states

3-spin system

##

 $C_7$ 

 $C_8$ 





x



Operates on entire system simultaneously

> Quantum Parallelism



3-spin system









### **Classical Gates**

GATE	CIRCUIT REPRESENTATION	TRUTH TABLE				<i>n</i>	
NOT The output is 1 when the input is 0 and 0 when the input is 1.	->>-	InputOutput0110		GATE	CIRCUIT REPRESENTATION	TRUTH TABLE	
AND The output is 1 only when both inputs are 1, otherwise the output is 0.	=D-	Input         Output           0         0           0         1           0         0           1         0           1         1	NOT	The output is 1 when the input is 0 and 0 when the input is 1.	->	Input         Output           0         1           1         0	<u>t</u>
OR The output is 0 only when both inputs are 0, otherwise the output is 1.	=D-	Input         Output           0         0           0         1           1         0           1         1				Input Output	t
NAND The output is 0 only when both inputs are 1, otherwise the output is 0.	⊐D⊷	Input         Output           0         0         1           0         1         1           1         0         1           1         1         0	AND	The output is 1 only when both inputs are 1, othe	-D-	$\begin{array}{c c} \hline mpar} \\ \hline 0 \\ 0 \\ 0 \\ 1 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\ \end{array} \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\ \end{array} \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\ \end{array} \\ \end{array} \\ \end{array} \\ \begin{array}{c} \hline 0 \\ \end{array} \\$	-
NOR The output is 1 only when both inputs are 0, otherwise the output is 0.	⊐⊅∽	Input         Output           0         0         1           0         1         0           1         0         0           1         1         0		output is C • Univ	versal gate se	ts for Boolean I	ogic
XOR The output is 1 only when the two inputs have different value, otherwise the output is 0.	⊐D-	Input         Output           0         0           0         1           1         0           1         1           0         0	OR	The outpur when both – E. are 0, othe	g., NOT, AND g., NOR		
XNOR The output is 1 only when the two inputs have the same value, otherwise the output is 0.		Input         Output           0         0         1           0         1         0           1         0         0           1         1         1		output is 1. – Ai – Re	nd many more (not equires at least on	: unique) e two-bit gate	

### Single-Qubit Quantum Gates



GATE	CIRCUIT REPRESENTATION	MATRIX REPRESENTATION	TRUTH TABLE	BLOCH SPHERE
I Identity-gate: no rotation is performed.	— <u>I</u> —	$I = \left(\begin{array}{cc} 1 & 0 \\ 0 & 1 \end{array}\right)$	$\frac{\text{Input}}{ 0\rangle}  \frac{\text{Output}}{ 0\rangle} \\  1\rangle   1\rangle$	z x y
X gate: rotates the qubit state by $\pi$ radians (180°) about the x-axis.	— <u>X</u> —	$X = \left(\begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array}\right)$	$\frac{\text{Input}}{ 0\rangle}  \frac{\text{Output}}{ 1\rangle} \\  1\rangle   0\rangle$	z 180° y
Y gate: rotates the qubit state by $\pi$ radians (180°) about the y-axis.	— <u>Y</u> —	$Y = \left(\begin{array}{cc} 0 & -i \\ i & 0 \end{array}\right)$	$\frac{\text{Input}}{ 0\rangle}  \frac{\text{Output}}{i  1\rangle} \\  1\rangle  -i  0\rangle$	

# Two-Qubit Quantum Gates



GATE	CIRCUIT REPRESENTATION	MATRIX REPRESENTATION	TRUTH TABLE
Controlled-NOT gate: apply an X-gate to the target qubit if the control qubit is in state  1)		$CNOT = \left(\begin{array}{rrrr} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{array}\right)$	$\begin{array}{c c} Input \\ \hline 100\rangle & 000\rangle \\ \hline 101\rangle & 101\rangle \\ \hline 110\rangle & 111\rangle \\ \hline 111\rangle & 110\rangle \end{array}$
Controlled-phase gate: apply a Z-gate to the target qubit if the control qubit is in state  1)		$cZ = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$	$\begin{array}{c c} \mbox{Input} & \mbox{Output} \\ \hline \mbox{I00} & \mbox{I00} \\ \hline \mbox{I01} & \mbox{I01} \\ \hline \mbox{I10} & \mbox{I10} \\ \hline \mbox{I11} & \mbox{-}\mbox{I11} \\ \end{array}$

- Universal gate sets for quantum logic
  - E.g., H, S, T, CNOT
  - And many more (not unique)
  - Requires at least one two-qubit entangling gate

### Single-Qubit Gate Example



#### Single-Qubit Gate Example



#### Microwave Pulse Control

#### **Gate Sequence**





#### **Application**



Control applied via capacitive or inductive coupling of a microwave pulse to the qubit

#### Microwave Pulse Control

#### X and Y Rotations on the Bloch Sphere

**I-Q Mixing** 





I: in-phase  $(0^{\circ}) \rightarrow x$ Qxis Q: quadrature  $(90^{\circ}) \rightarrow y$  axis

## Two-Qubit Gate Example



For example:

$$|\psi_{\text{out}}\rangle \propto |0\rangle_{x} |0\rangle_{y} + |1\rangle_{x} |1\rangle_{y} \neq (\cdots)_{x} (\cdots)_{y}$$

 $|\psi\rangle \propto (|0\rangle + |1\rangle) |0\rangle$ 

Results in an *entangled state* (cannot be factored)

Universal gate-model quantum computation is achievable with a small set of single and two-qubit gates.

# Quantum Algorithm



### Paths to Applications



M. Kjaergaard, WDO, et al., Annual Reviews of CMP 11, 369-395 (2020)

#### Commercial Quantum Advantage



Small region where useful quantum algorithms exist (as we know them today)

# Types of Quantum Advantage



### **Exponential Growth**

#### Exponential Growth: Doubling Pennies Every Day for 1 Month





 $2^0 = 1$  penny

 $2^{1} = 2$  pennies  $2^{2} = 4$  pennies  $2^{3} = 8$  pennies

# After 31 days, would you take the pennies or \$10M?

### **Exponential Growth**

#### Exponential Growth: Doubling Pennies Every Day for 1 Month





2<sup>1</sup> = 2 pennies 2<sup>2</sup> = 4 pennies 2<sup>3</sup> = 8 pennies

#### 2<sup>31</sup> = 2,147,483,648 pennies > \$21M !!

#### **Exponential Power**

• Simulating quantum computers (QCs) on classical computers



#### **Exponential Power**

• Simulating quantum computers (QCs) on classical computers

Qubits	Size of simulator
30	laptop
50	supercomputer
### **Exponential Power**

• Simulating quantum computers (QCs) on classical computers

Qubits	Size of simulator
30	laptop
50	supercomputer
80	all computers on Earth

### **Exponential Power**

• Simulating quantum computers (QCs) on classical computers

Qubits	Size of simulator
30	laptop
50	supercomputer
80	all computers on Earth
160	all Si atoms in Earth

### **Exponential Power**

• Simulating quantum computers (QCs) on classical computers

Qubits	Size of simulator
30	laptop
50	supercomputer
80	all computers on Earth
160	all Si atoms in Earth
300	> all atoms than in known universe

# Digital Quantum Algorithms

Algorithm	Classical Time	Quantum Time	Speedup	Limitation	
Simulation <sup>1</sup> (quantum chemistry)	2 <sup>N</sup> (for N atoms)	Nc	Exp. in space, polynomial in time	Mapping problem to qubits	
<b>Factoring</b> <sup>2</sup> (+ related number theoretic)	2 <sup>N</sup> (for N digits)	N <sup>3</sup>	Exponential	Classical runtime limit unproven	
Linear systems <sup>3</sup> (Ax=b)	2 <sup>N</sup> (for N digits)	~N	Exponential	Strict conditions, e.g. sparse matrix	
<b>Optimization</b> <sup>4</sup>	2 <sup>N</sup>	?	?	Empirical	
Search <sup>5</sup> (unsorted / unstructured data)	Ν	$\sqrt{N}$	<b>Polynomial</b> $(\sqrt{N})$	Data loading	
Anand Natarjan	Ike ChuangSeth Lloyd	1,3 Peter Shor <sup>2</sup> Ara	$am Harrow^{3}$	Michael Sipser <sup>4</sup>	

### Minimum requirements for the physical implementation of a quantum computer

- D1: Robust, reproducible qubit technology
- **D2: Initialization**
- **D3: Measurement**
- D4: Universal set of gates
- **D5: Coherence & fidelity**

#### **Dedicated Superconducting Qubit Fab**

- 200-mm wafers & 50-mm wafers
- Qubits and classical digital electronics
- Deposition, dry etch, PECVD, CMP
- Unique facility worldwide

#### **Custom Plassys Evaporator**



### **Electron Beam Lithography**





MIT-LL Raith EBPG5200 routinely patterns <150 nm Josephson junctions

### Veeco Gen-200 MBE



### Minimum requirements for the physical implementation of a quantum computer

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#### **Fabrication Process Monitoring**

- Data-driven process development
- >1000-10,000 test structures (50-200 mm wafers)
- JJs, lines, combs & snakes, contacts, crossovers, chains, ...
- Automated testing and analysis

#### **Room-Temp Probe Station**



#### **Cross-Wafer Variation Maps** 4.18 4.33 4.43 4.28 3.98 4.49 4.14 4.06 4.57 4.55 4.34 4.16 4.09 4.52 4.42 4.21 4.58 3.9 4.58 4.46 4.23 J = 4.2765 atd = 0.21131

#### S. Tolpygo, ..., WDO, IEEE Appl. Supercond. (2014, 2015); K.K. Berggren et al., ibid (1999)



WDO, et al., unpublished (2006)

### Minimum requirements for the physical implementation of a quantum computer

D1: Robust, reproducible qubit technology

#### **D2: Initialization**

**D3: Measurement** 

D4: Universal set of gates

#### **D5: Coherence & fidelity**

### High-fidelity (99.9%) state initialization

• Microwave cooling (active)

55

60

65

50

- Cryogenic engineering (passive)
- Active measurement-based feedback

#### **Microwave Cooling**

WDO et al., Science 310, 1653 (2005); Science 310, 1589 (2006); Nature 455, 51 (2008)



#### Cryogenic Engineering

#### X. Jin, ..., WDO, PRL 114, 240501 (2015)

pexp: Maxwell-Boltzmann (Eq. 1)

Ples : Maxwell-Boltzmann

2-point averaging

٥.

15

20

residual population ~ 0.1%

25

30

35

Bath Temperature (mK)

40

45

Excited State Population



A. Greene, ..., WDO, APS MM (2018)



### Minimum requirements for the physical implementation of a quantum computer

D1: Robust, reproducible qubit technology

- D2: Initialization
- D3: Measurement
- D4: Universal set of gates

### D5: Coherence & fidelity

### High-fidelity (99%) measurement

- Control electronics and software
- Syndrome measurement and feedback
- Error detection and correction

Gustavsson, Krantz, Hover, and WDO

Control electronics, software, and quantum-limited amplifiers high-fidelity measurement of error syndromes





Oubit

Contro



**D**Labber

Macklin, WDO, et al., Science (2015)



### Minimum requirements for the physical implementation of a quantum computer

- D1: Robust, reproducible qubit technology
- D2: Initialization
- D3: Measurement
- D4: Universal set of gates
- D5: Coherence & fidelity











### Coherence Time and Gate Time



**Gate time t**<sub>gate</sub>: Time required for a single gate operation

**Figure of Merit** \* : # of gates per coherence time =  $t_{coh}/t_{gate}$ 

(\* Rigorous metric: gate & readout fidelity)

Long coherence times are not sufficient, it's the number of gates before an error

## **Qubit Modalities**



# **Qubit Modalities**



**MIT Campus** 

#### **MIT Lincoln Lab**





Rajeev Ram Ike Chuang Physics, EECS

John Chiaverini LL, RLE



Will Oliver EECS, Phys., LL





EECS

Terry Orlando Jamie Kerman

LL

and large teams at MIT & LL

EECS

# **Qubit Modalities**



# Artificial Atom: Superconducting Qubits

### Qubit: superconducting circuit

Phase, flux, or charge

- □ Coherence times: ~ 100 us
- □ Fidelity and operation times

1 QB:	99.99% in 10 ns
2 QB:	99.9% in 40 ns
_	00.00/ : 000

- Readout: 99.0% in 200 ns
   Clock rate: ~ 25 MHz
- □ Largest algorithm: 53 qubits
- □ Companies:
  - AWS, Google, IBM, QCI, Rigetti, ...
  - Annealing: D-Wave

#### **Electrical Circuit -- Anharmonic Oscillator**





# Atomic State: Trapped Ion Qubits

- Qubit: energy levels of an ionized atom
  - Ca+, Sr+, Be+
  - Optical or microwave transitions
- □ Coherence times: 10 s
- □ Fidelity and operation times
  - 1 QB: 99.999% in 5 us
    2 QB: 99.900% in 50 us
    Readout: 99.990% in 30 us
- □ Clock rate: ~ 20 kHz
- □ Largest algorithm: 30 qubits
- Companies: Honeywell, Ion-Q, AQT, Universal Quantum, ...





## Atomic State: Neutral Atoms

- Qubit: energy levels of a neutral atom
  - Rb, Cs, Ho trapped in an optical lattice
  - Optical and microwave fields
- □ Coherence times: 1 s
- □ Fidelity and operation times
  - 1 QB: 99% in 3 us
     2 QB: >99% in 300 us
     Readout: 99.90% in >3 milliseconds
- □ Clock rate: 10 kHz
- □ Largest lattices: 100-300 qubits
- Companies: Atom Computing, ColdQuanta, Pasqual, QuEra



## Electron Spin: SiGe Quantum Dots

### Qubit: electron spin

- Quantum dots in SiGe 2DEGs
- RF and baseband pulsing
- Double-dot, triple-dot, CMOS dot
- Coherence times: 400 us

### Fidelity and operation times

- 1 QB: 99.5% in 100 ns
- 2 QB: >99% in 200 ns
- Readout: 99% in 1 us
- □ Clock rate: 5 MHz
- Companies: HRL, Intel

#### SiGe Quantum Dots



#### **Energy Levels**



Nature 479, 345 (2011)

# Electron Spin: Phosphorus-Doped Silicon

- Qubit: electron spin (nuclear spins)
  - Phosphorus donor in silicon
  - Microwave pulses
- Coherence times: 100 ms (1 s)

### Fidelity and operation times

- 1 QB: 99.5% in 200 ns (99.99% in 100 us)
  - 2 QB: ~ 90% in 1-100 ns
- Readout: 95.0% in 1 ms (99.9% in 50 ms)
- Clock rate: TBD
- Companies: SQC (Silicon Quantum Computing)

#### **Phosphorous-Doped Silicon**





# Electron and Nuclear Spins: NV Centers

#### Qubit: electron or nuclear spin

- Nitrogen vacancy electron (NV-)
- Nitrogen or carbon nuclear spins
- Other defects may be used

### Coherence times: 20 ms

Fidelity and operation times

1 QB:	99.5% i	in 10 us
2 QB:	>90%	in 25 ι

>90% in 25 us

#### 94.0% in 50 us Readout:

### Clock rate: 40 kHz

Companies: N/A (mostly sensing applications)

#### **Diamond with Nitrogen Vacancy**



(note: redraw and have all carbon atoms be blue with a nuclear spin. Do not label C1...C4, just put an "n" inside one. Put an "e" inside the electron instead of e-. Label B rather than Bz

#### **Energy Levels**



# Benchmarking Methods

#### **Randomized Benchmarking**



### Pros

- Simple and efficient procedure to obtain fidelity
- Current 'gold standard'

### Cons

• Time dependent errors may alter decay curve

### **Quantum Process Tomography**



### Pros

 Exact reconstruction of any quantum process

### Cons

- Exponential resource requirement (3 qubits is the borderline)
- Cannot separate gate
   errors from SPAM errors

### Gate Fidelities

### Single-Qubit Gate Fidelity > 0.999

Two-Qubit Gate Fidelity > 0.995



Y. Sung et al. arXiv:2011.01261 (2020): Experiment

## Quantum Advantage Demonstrations



TABLE I. The runtime of tensor network algorithm for different circuits on Summit. The classical simulation consumption estimation of the random quantum circuit sampling experiment on the Sycamore, *Zuchongzhi* 2.0, and *Zuchongzhi* 2.1 processors are provided. FPOs is the abbreviation for the number of floating point operations, QPU is the abbreviation for quantum processing unit.

Processor	Circuit	Fidelity	# of bitstrings	FPOs (a perfect sample)	FPOs (circuit)	Runtime on	Runtime on	ClassicalRuntime QauntumRuntime
Sycamore [8]	53-qubit 20-cycle	0.224%	$3.0  imes 10^6$	$\frac{(a \text{ perfect sample)}}{1.63 \times 10^{18}}$	$1.10 \times 10^{22}$	15.9 days	600s	$2.29 \times 10^{3}$
Zuchongzhi 2.0 [11]	56-qubit 20-cycle	0.0662%	$1.9  imes 10^7$	$1.65 imes10^{20}$	$2.08\times 10^{24}$	8.2 years	1.2h	$6.02  imes 10^4$
Zuchongzhi 2.1	60-qubit 22-cycle	0.0758%	$1.5  imes 10^7$	$1.06  imes 10^{22}$	$1.21\times 10^{26}$	$4.8 \times 10^2$ years	1h	$4.21 \times 10^6$
Zuchongzhi 2.1	60-qubit 24-cycle	0.0366%	$7.0  imes 10^7$	$4.68  imes 10^{23}$	$1.2\times 10^{28}$	$4.8 \times 10^4$ years	4.2h	$9.93  imes 10^7$



The Google Quantum AI team demonstrated

a calculation in ~200s with one chip, 53 superconducting qubits, drawing around 100 kW of power

On the Summit supercomputer (Oak Ridge National Laboratory),

it would take several days, with all 40,000 CPUs & GPUs, 10<sup>17</sup> transistors & memory, and 100's MW of power

Google AI, Nature 505, 574 (2019); USTC, arXiv:2109.03494 (2021)

## Architectural Layers of a QIP



## Architectural Layers of a QIP

#### **Layered Architecture**



N.C. Jones PRX 2, 031007 (2012)

## Architectural Layers of a QIP

Engineered Error Mitigation: Dynamical Decoupling

#### Eg. Lacrosse Cradling





### Lacrosse in the Presence of Noise





## Dynamical Decoupling from Running "Noise"





### "Active Error Correction" in Lacrosse





# **Coherence** Times



# **Qubit Dephasing and Filter Function**



Engineered filter function depends on pulse sequence and windows the PSD  $S_{\lambda}(\omega)$ 

J. Bylander, ..., WDO, Nature Physics (2011), Martinis et al., PRB (2003), Ithier et al., PRB (2005); Yoshihara et al., PRL (2006), Cywinski et al. PRB (2008)

### **Dynamical Decoupling:** Noise Shaping Filters



### **Dynamical Decoupling:** Noise Shaping Filters with 1 π-pulse



### **Dynamical Decoupling:** Noise Shaping Filters with 2 π-pulses



### **Dynamical Decoupling:** Noise Shaping Filters with $N \pi$ -pulses



#### Carr – Purcell (– Meiboom – Gill) Sequence

 $\frac{\tau}{N}$ 

10

4

Frequency, f (MHz)

 $\frac{\tau}{N}$ 

2N

 $\tau = 1 \ \mu s$ 

S ~ 1/f

6

CP(MG) UDD

8

### Noise Spectroscopy

#### Qubit Noise Spectroscopy Filter Engineering & Optimal Control



#### Y. Sung, ..., WDO, Nature Communications 10, 3715 (2019) F. Yan, ..., WDO, Nature Communications 7, 12964 (2016) F. Yan , ..., WDO, Nature Communications 4, 2337 (2013)

#### Carr – Purcell (– Meiboom – Gill) Sequence



## Dispelling Myths About QC

- Quantum computers will not replace classical computers
- Quantum computers will not break encryption soon
  - RSA 2048-bit keys: around 4000 error corrected qubits
  - Bitcoin encryption: around 2300 error corrected qubits
- However, one should not wait until a quantum computer can break RSA to switch to post-quantum encryption


## Quantum Engineering



Quantum Engineering is the bridge connecting science, mathematics, and classical engineering

# Randomized Benchmarking

Single-qubit randomized benchmarking



J. Emerson *et al. J. Opt. B* 7, S347 (2005)
 E. Knill *et al. Phys. Rev. A*. 77, 012307 (2008)
 E. Megesan *et al. Phys. Rev. Lett.* 106, 180504 (2011)

- Goal: estimate the average error rates of quantum gates.
- (Clifford-based) Randomized Benchmarking [1,2,3]
  - Initialize qubits at the ground state.
  - Apply *m* randomly chosen Clifford gates ( $C_1$ ,  $C_2$ , ...,  $C_m$ ).
  - At the end, apply the inverse gate s.t. the entire operation = Identity.
  - Measure the survival probability of the ground state (= "sequence fidelity"  $F_{seq}$ ).
    - ✓ In the absence of error  $\rightarrow$   $F_{seq} = 1$ .
    - ✓ In the presence of error  $\rightarrow$   $F_{seq} < 1$ .

Twirling over Cliffords  $\rightarrow$  Depolarization of the gate error [1,2,3]

$$ho 
ightarrow p
ho + rac{(1-p)}{2^n}I$$
 (*n*: # of qubits)

- $F_{seq}$  will decay exponentially as  $F_{seq} = Ap^m + B$ .
- The average error rate per Clifford  $r_{\text{Cliiford}}$  is related to p as

$$r_{\text{Clifford}} = (1-p) \times \frac{2^n - 1}{2^n}$$

#### Randomized Benchmarking

Single-qubit randomized benchmarking





- Goal: estimate the average error rates of quantum gates.
- (Clifford-based) Randomized Benchmarking [1,2,3]
  - Initialize qubits at the ground state.
  - Apply *m* randomly chosen Clifford gates ( $C_1$ ,  $C_2$ , ...,  $C_m$ ).
  - At the end, apply the inverse gate s.t. the entire operation = Identity.
  - Measure the survival probability of the ground state (= "sequence fidelity" F<sub>seq</sub>).
    - ✓ In the absence of error  $\rightarrow$   $F_{seq} = 1$ .
    - In the presence of error  $\rightarrow F_{seq} < 1$ .

Measurement of the avg. error rate per 1QB Clifford

- Fit  $F_{seq}$  with exponential  $(f(x) = Ae^{Bx} + C)$ .
- Extract depolarizing rate p, where  $p = e^B$ .
- *A*, *C*: absorbs the SPAM error.
- The average error rate per 1QB Clifford gate r,

$$r_{\text{Clifford}} = (1-p) \times \frac{2^n - 1}{2^n} = \frac{1-p}{2}$$

(n: # of qubits)

#### Interleaved Randomized Benchmarking

- **Reference** (1QB) randomized benchmarking ( $F_{seq, ref}$ )





- Interleaved (1QB) randomized benchmarking ( $F_{seq, int}$ )



- Interleaved Randomized Benchmarking [1,2]
  - Interleave gate of interest G at every Clifford ( $C_1, C_2, ..., C_m$ ).
  - Compare it to the reference RB to extract the error rate of *g*.
  - $F_{\text{seq, ref}} = Ap_{\text{ref}}^{m} + B$  (reference curve)
  - $F_{\text{seq, int}} = A' (p_{\text{ref}} p_g)^m + B' \equiv A' p_{\text{int}}^m + B' \text{ (interleaved curve)}$
- The average error rate per interleaved gate  $r_{
  m int}$ ,

$$r_{\text{int}} = \left(1 - \frac{p_g}{p_g}\right) \times \frac{2^n - 1}{2^n} = \left(1 - \frac{p_{\text{int}}}{p_{\text{ref}}}\right) \times \frac{2^n - 1}{2^n}$$

[1] E. Megesan *et al. Phys. Rev. Lett.* **109**, 080505 (2012)
[2] A. D. Corcoles *et al. Phys. Rev. A.* **87**, 030301 (2013)

# Two-Qubit Randomized Benchmarking

- **Reference** (2QB) randomized benchmarking ( $F_{seq, ref}$ )



- Interleaved (2QB) randomized benchmarking ( $F_{seq, int}$ )



- Measurement of the avg. error rate per interleaved gate
  - Sequence fidelity  $F_{seq}$  = the survival probability of  $|00\rangle$ .
  - $F_{\text{seq, ref}} = Ap_{\text{ref}}^{m} + B \text{ (reference curve)}$
  - $F_{\text{seq, int}} = A' p_{\text{int}}^m + B' \text{ (interleaved curve)}$
  - The average error rate per interleaved gate (CZ)  $r_{\rm int}$ ,

$$r_{\rm int} = \left(1 - \frac{p_{\rm int}}{p_{\rm ref}}\right) \times \frac{2^n - 1}{2^n} = \left(1 - \frac{p_{\rm int}}{p_{\rm ref}}\right) \times \frac{3}{4}$$

• Avg. CZ fidelity  $F_{CZ} = 1 - r_{int}$