# Virtual Machines in Distributed Environments

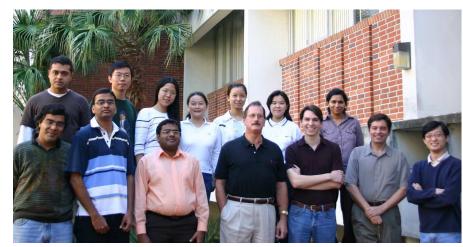
Maurício Tsugawa on behalf of José A. B. Fortes

Advanced Computing and Information System Laboratory University of Florida, Gainesville



# **ACIS Lab history + statistics**

- Founded 8/10/01
- 20+ people
- 5 "associated" faculty
  - ECE, CISE, MSE



- Also 10+ "associated" faculty from Purdue, CMU, U.
  Colorado, Northwestern U., U. Mass., NCSU ...
- 7 M dollars in funding, 5 M dollars in equipment
  - Approximately 1.5 M of subcontracts
  - NSF, NASA, ARO, IBM, Intel, SRC, Cisco, Cyberguard ...
- Computer infrastructure/lab
  - 300+ CPUs, .5 Tflops, 9 TBytes, Gigabit connections
  - Access to CPUs at Purdue, Northwestern, Stevens I.



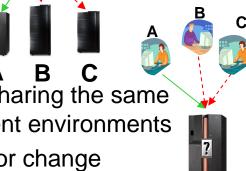
#### **Outline**

- What's in a talk title
  - Environment as a container for app execution
  - Distributed "a la" power grid
  - Virtualization for creation and coexistence of different environments in physical resources
- A Grid-building recipe
- Words are easy, let's build it: In-VIGO
  - Architecture, Deployments and Futures
  - Virtual Machines, Data, Networks and Applications
  - Turning Applications into Grid Services
- Conclusions



# **Resource sharing**

- Traditional computing/data center solutions:
  - Multitask/multiuser operating systems, user accounts, file systems ...
  - Always available but static configurations
    - Sharing possible if apps run on similar execution environments
  - Centralized administration
    - Tight control on security, availability, users, updates, etc
- Distributed Grid/datacenter requirements
  - Multiple administrative domains
    - Different policies and practices at each domain
    - Many environments possible
    - Dynamic availability
  - Must run all kinds of applications
    - Application user will neither trust unknown users sharing the same resource nor redevelop application to run in different environments
    - Resource owner will neither trust arbitrary users nor change environment for others' applications



### "Classic" Virtual Machine

- Copy of a real machine
  - "Any program run under the VM has an effect identical with that demonstrated if the program had been run in the original machine directly" 1
- Isolated from other virtual machines
  - "...transforms the single machine interface into the illusion of many"
- Efficient
  - "A statistically dominant subset of the virtual processor's instructions is executed directly by the real processor"
- Also known as a "system VM"



<sup>&</sup>lt;sup>1</sup> "Formal Requirements for Virtualizable Third-Generation Architectures", G. Popek and R. Goldberg, Communications of the ACM, 17(7), July 1974

<sup>&</sup>lt;sup>2</sup> "Survey of Virtual Machine Research", R. Goldberg, IEEE Computer, June 1974

### **Process vs. System VMs**

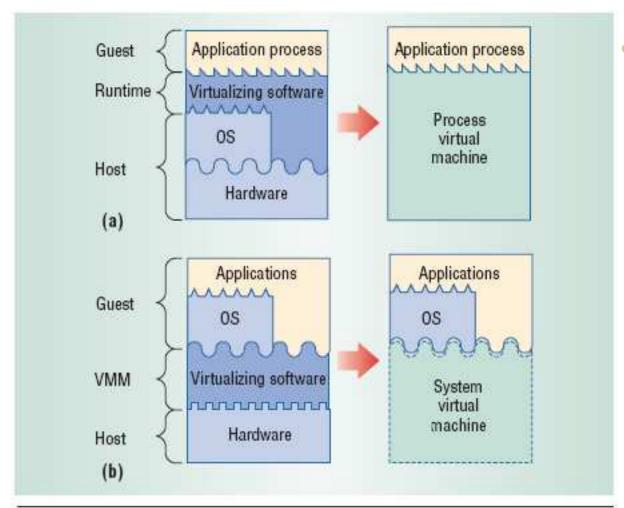
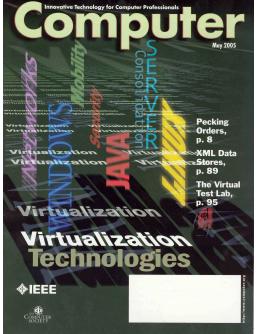


Figure 3. Process and system VMs. (a) In a process VM, virtualizing software translates a set of OS and user-level instructions composing one platform to those of another. (b) In a system VM, virtualizing software translates the ISA used by one hardware platform to that of another.

 In Smith and Nair's "The architecture of Virtual machines", Computer, May 2005

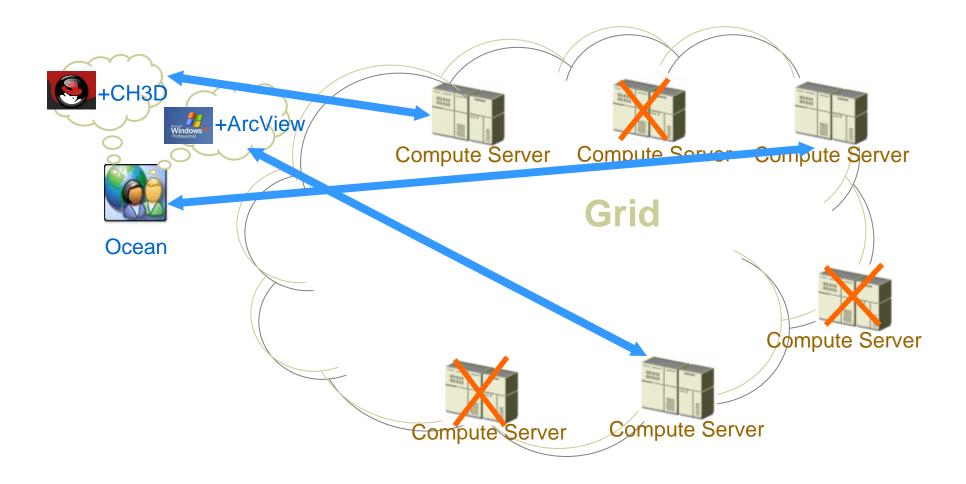


#### **Classic Virtual Machines**

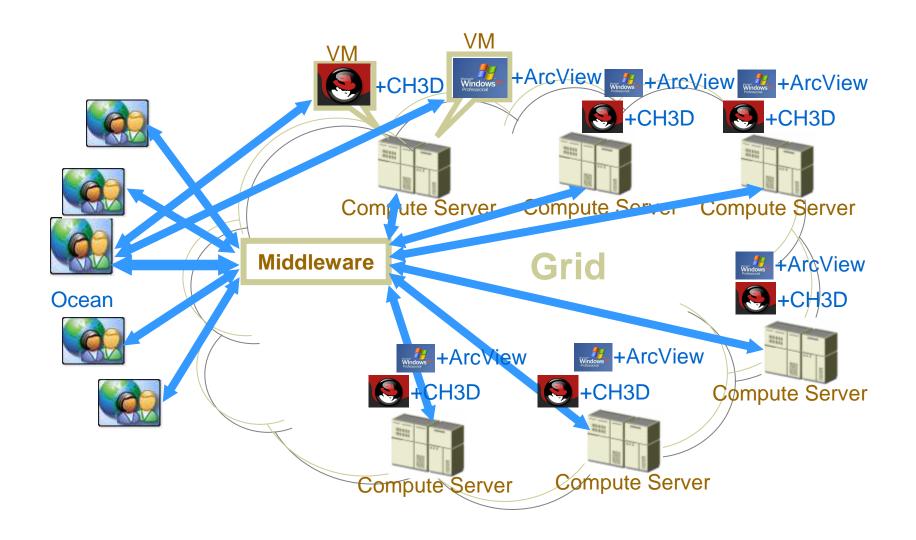
- Virtualization of instruction sets (ISAs)
  - Language-independent, binary-compatible (not JVM)
- 70's (IBM 360/370..) 00's (VMware, Microsoft Virtual Server/PC, z/VM, Xen, Power Hypervisor, Intel Vanderpool, AMD Pacifica ...)
- ISA+ OS + libraries + software = execution environment



### 1 user, 1 app, several environments



### Many users, 1 app, many environments



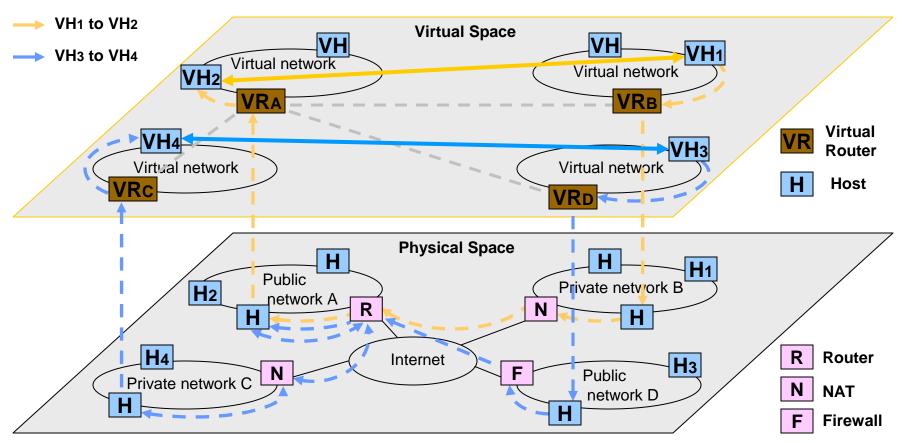


### Virtualization technology for grids

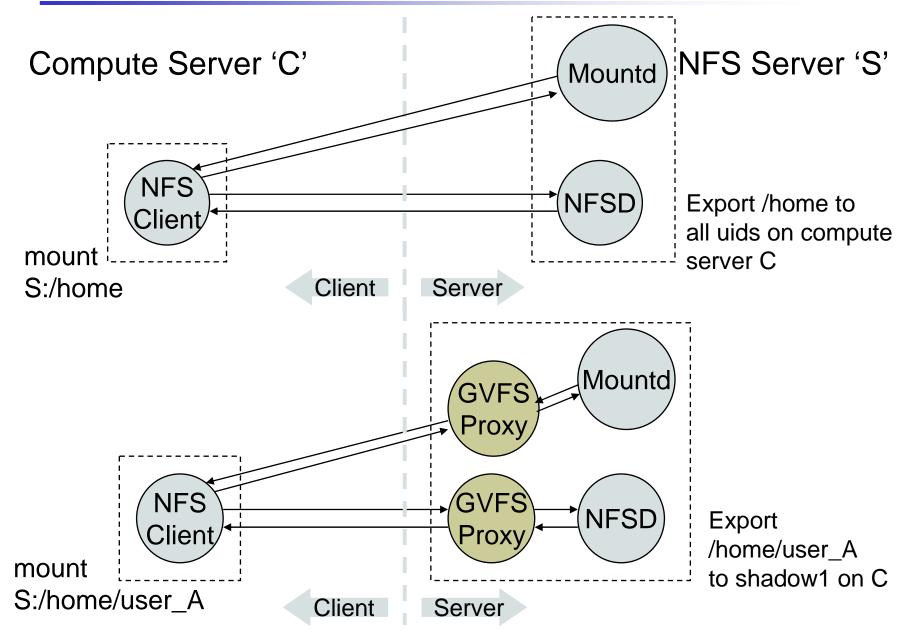
- Resource virtualization technology
  - Enables a resource to <u>simultaneously</u> appear as <u>multiple</u> resources with possibly <u>different</u> <u>functionalities</u>
    - Polymorphism, manifolding and multiplexing
- Virtual networks, data, applications, interfaces, peripherals, instruments ...
  - Emergent technologies

#### Virtual networks

- logical links:
  - multiple physical links, routing via native Internet routing
  - tunneling, virtual routers, switches, ...
  - partial to total isolation



#### **Data/file virtualization**

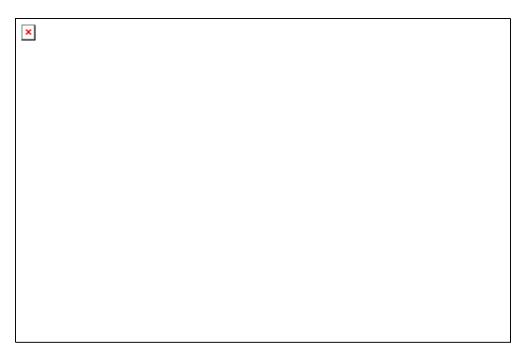


#### **Web services framework**

allows programs on a Service network to find each Description other, communicate **Service** and interoperate by Registry using standard protocols and languages **Publish** Find WSDL, UDDI WSDL, UDDI Service **Service** provider Service requestor Bind Service Description

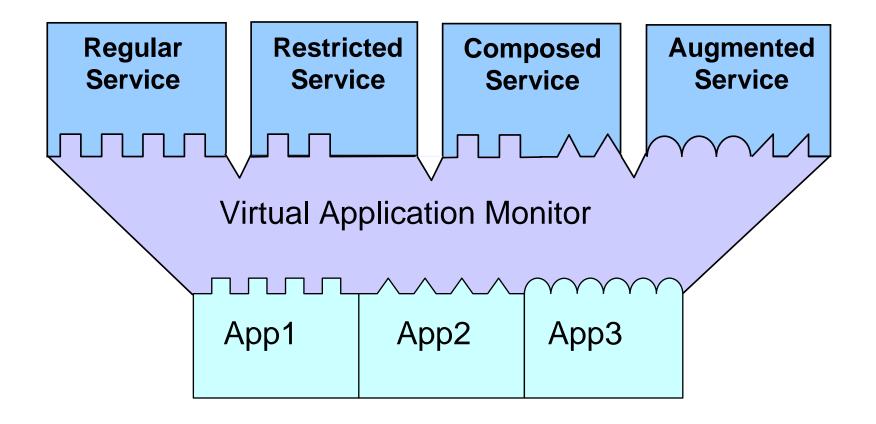
#### **Basic service description: interface definition**

 abstract or reusable service definition that can be instantiated and referenced by multiple service implementation definitions



 different implementations using the same application can be defined to reference different service definitions – a form of virtualization

### **Application virtualization**

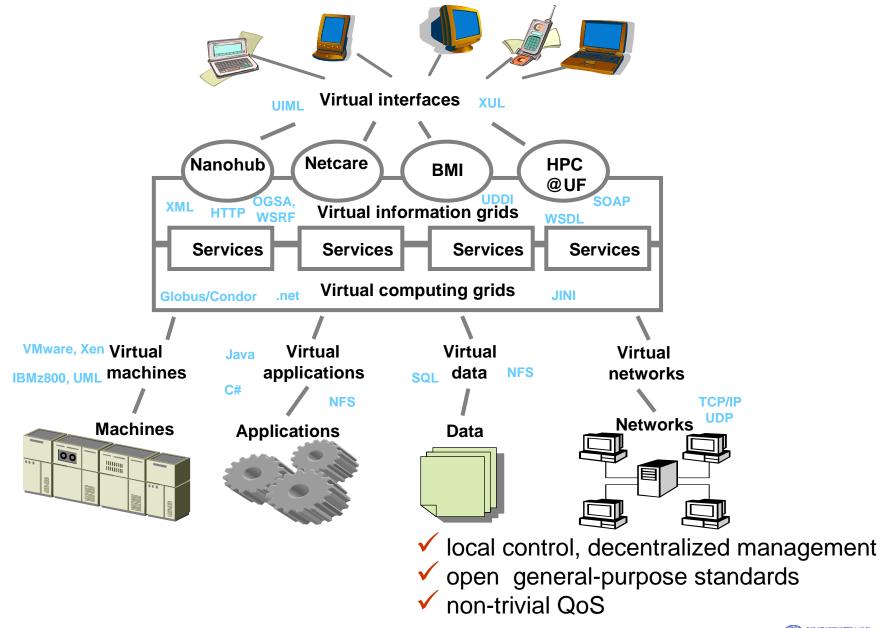


# A Grid-building recipe

- Virtualize to fit needed environments
- Use services to generate "virtuals"
- Aggregate and manage "virtuals"
- Repeat **000** as needed
- Net result:
  - users interact with virtual entities provided by services
  - middleware interacts with physical resources
- In-VIGO is a working proof-of-concept!



### The In-VIGO approach

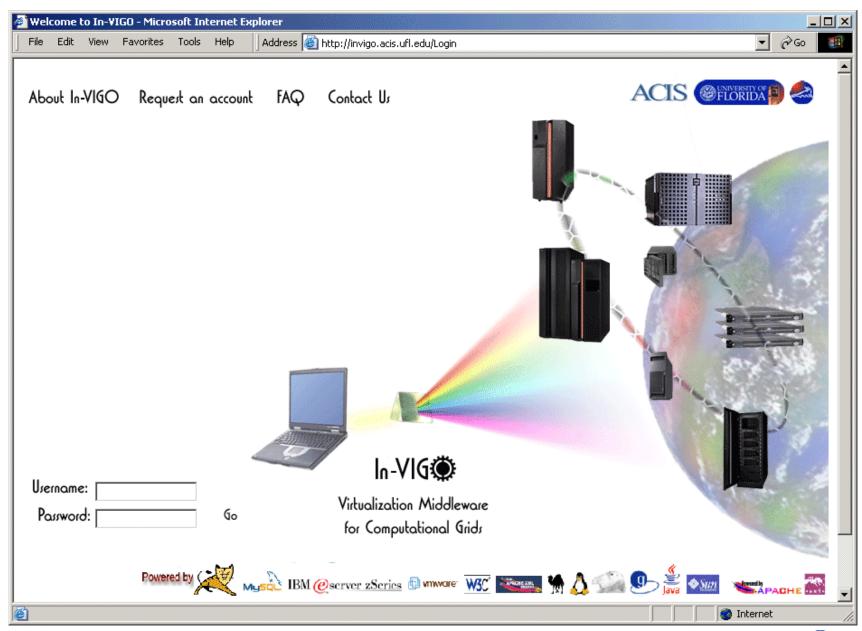


#### In-VIGO: a user's view

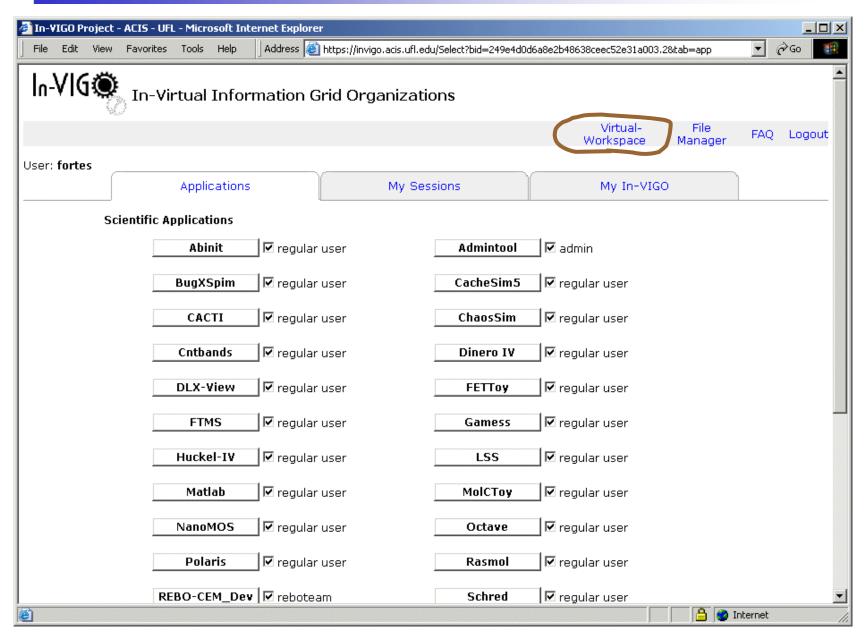
- Enables computational engineering and science <u>In-Virtual</u>
  Information <u>Grid Organizations</u>
- Motivations:
  - Hide complexity of dealing with cross-domain issues
    - From application developers
    - From end users
  - Provide secure execution environments
- Goals:
  - Application-centric: support unmodified applications
    - Sequential, parallel
    - Batch, interactive
    - Open-source, commercial
  - User-centric: support Grid-unaware users



# http://invigo.acis.ufl.edu

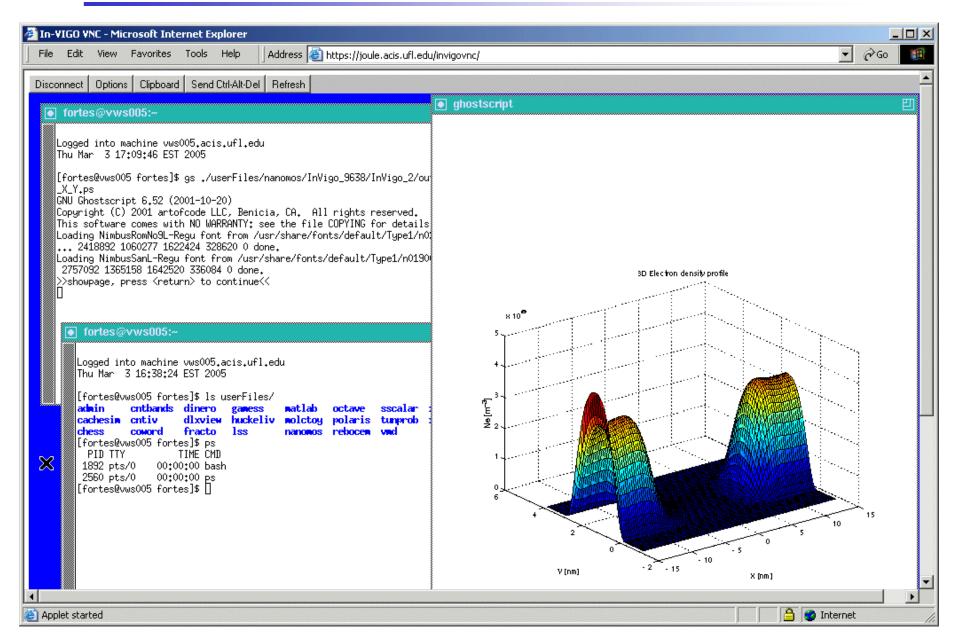


# The In-VIGO portal



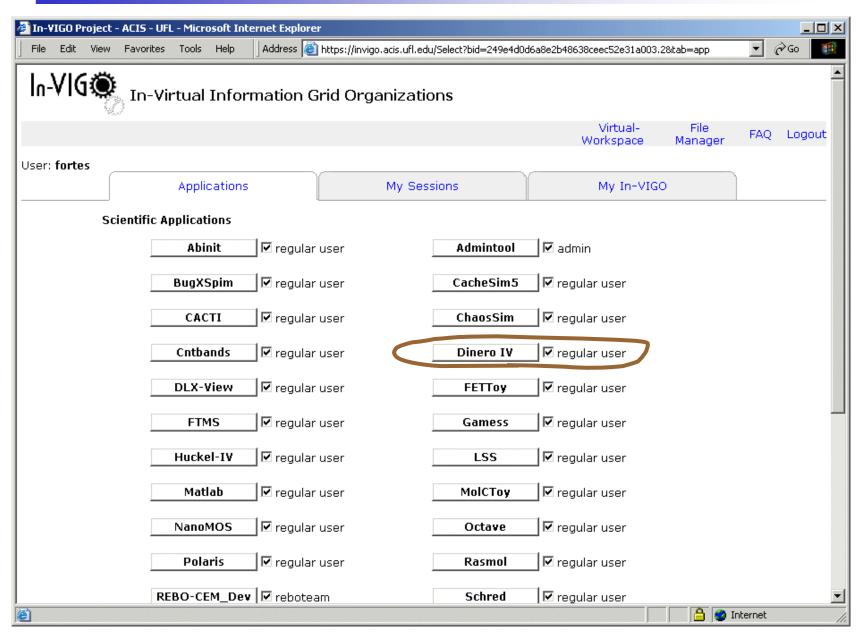


### Virtual workspace





# The In-VIGO portal



# Setting up ...

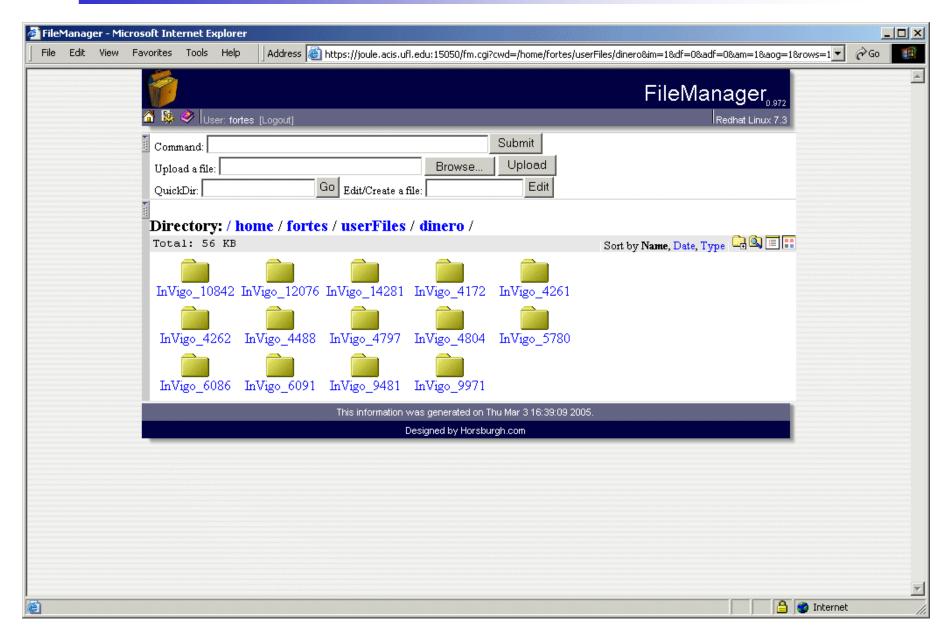


#### Interface and workflow

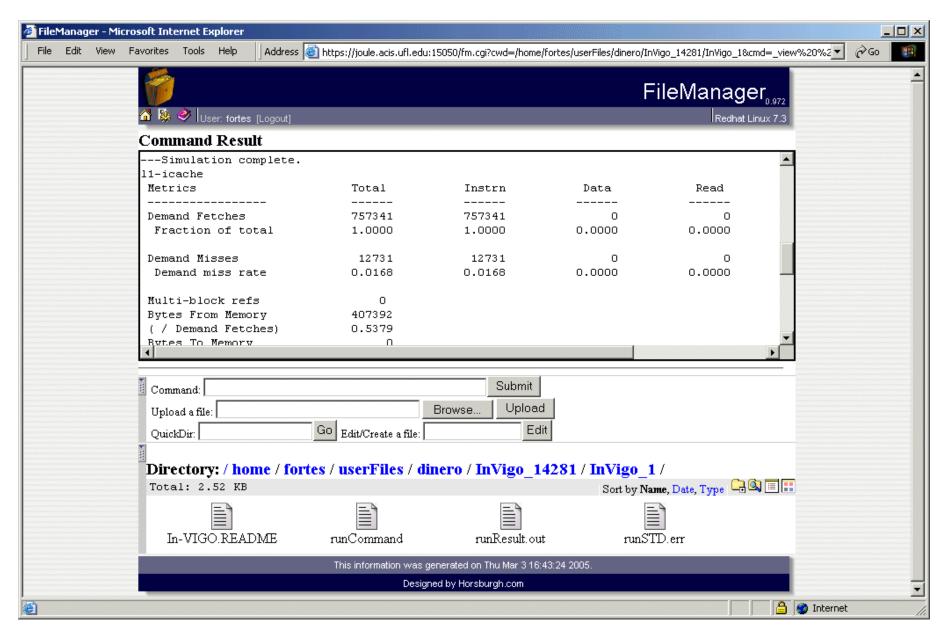




# File manager (1)

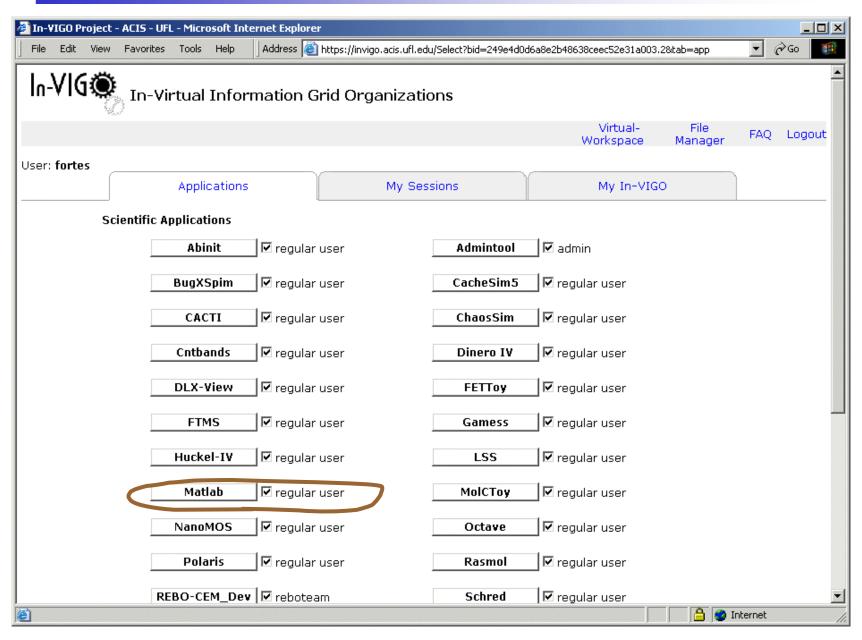


# File manager (2)

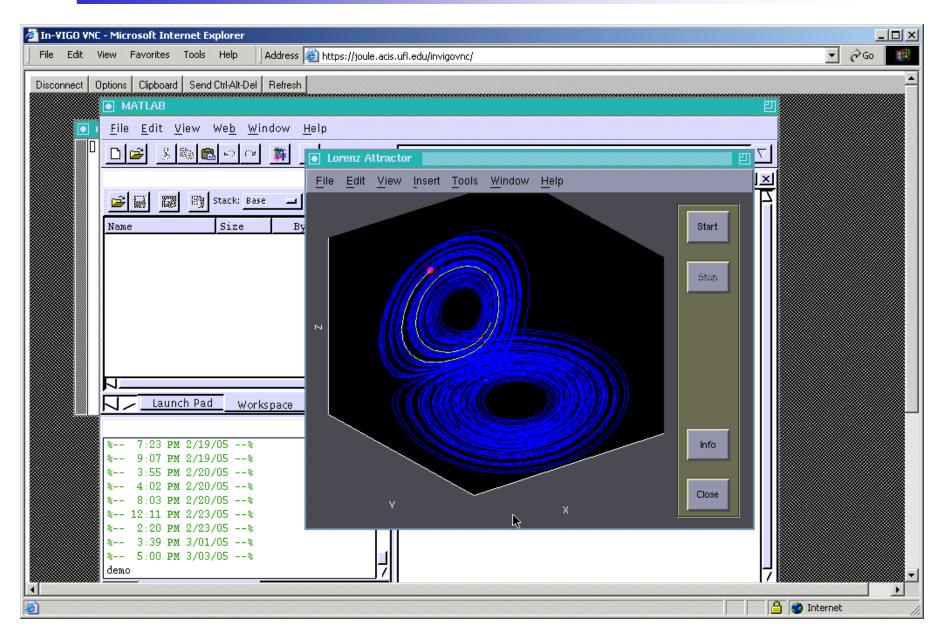




# The In-VIGO portal

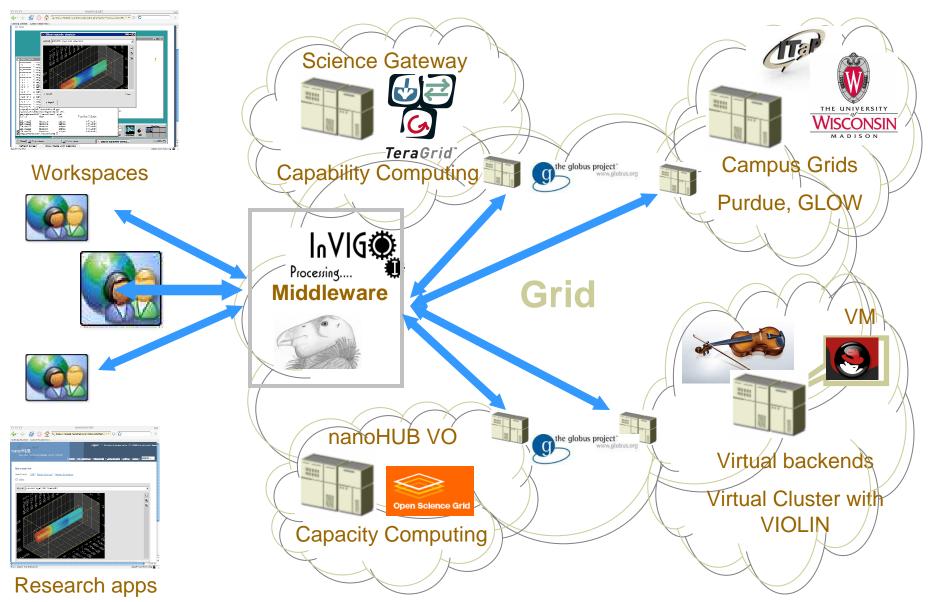


#### **Native interactive interface**





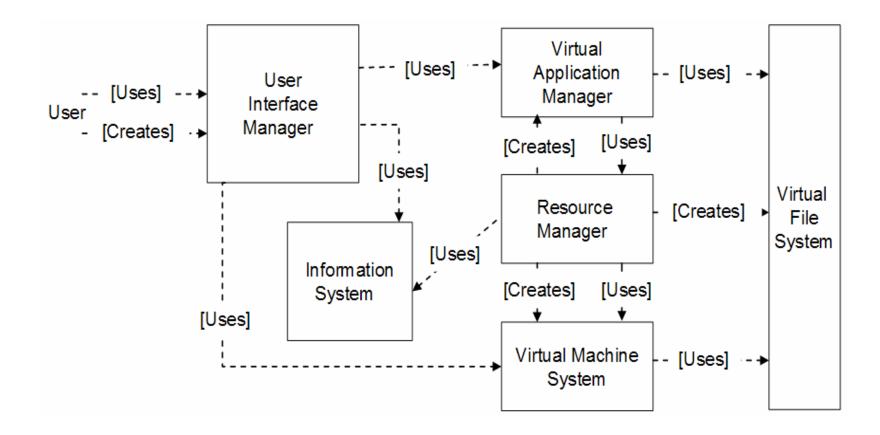
#### nanoHUB (current middleware infrastructure)





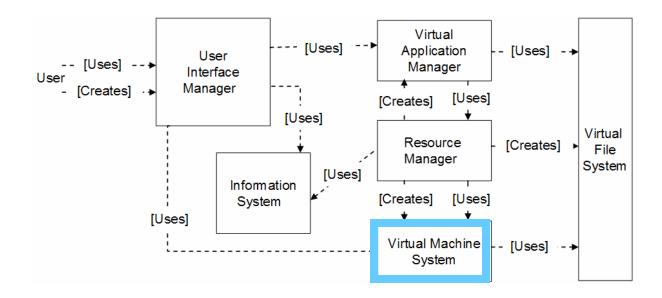
### In-VIGO 1.0 architecture diagram

- Deployed at UF/ACIS since Fall 2003
- nanoHUB: Summer 2005
- On-going deployments: SURA/SCOOP, UF/HPC



# **Virtual Machine System**

- Enables on-demand instantiation of whole-O/S VMs for virtual workspaces
- Has access to physical resources (host)
- Create, configure, query, destroy VMs
- In-VIGO users have access to virtual resource (guest)



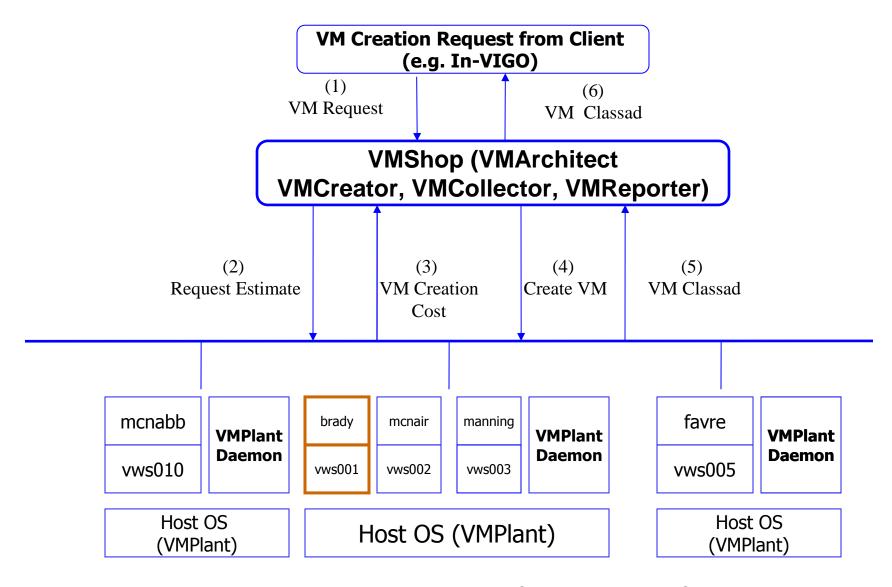
#### VM services

- Provide means to efficiently create/configure/destroy VMs,
  - generic across VM technologies
- Directed Acyclic Graph (DAG) model for defining application-centric VMs
- Cost-bidding model for choosing compute servers for VM instantiation

(SC 2004)



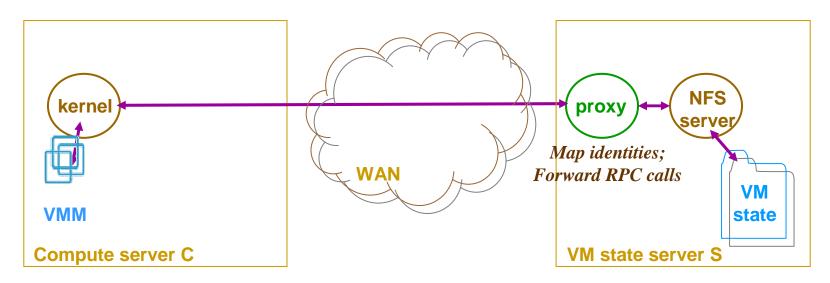
#### **Architectural Components of VM Service**



#### **Data access virtualization**

- Grid virtual file systems (GVFS)
  - On-demand setup, configuration and tear-down of distributed file systems
    - Unmodified applications access file-based data in the same manner they would in a local environment
    - Use and extend Network File Systems (NFS)
  - Multiple, independent file system sessions share one or more accounts in file servers
  - File system data is transferred on-demand, on a per-block basis

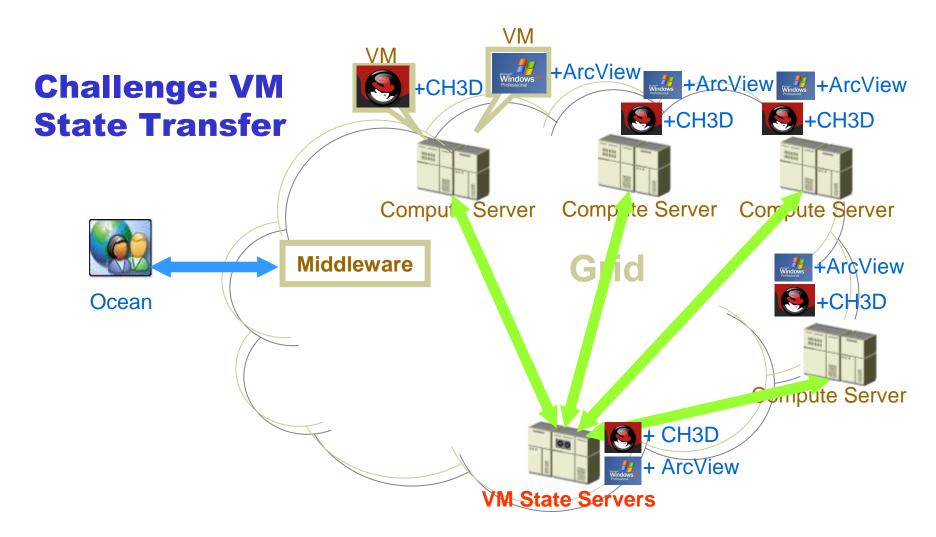
# **Grid Virtual File System (GVFS)**



- Logical user accounts [HCW'01] and virtual file system [HPDC'01]
  - Shadow account + file account, managed by middleware
  - NFS call forwarding via middle tier user-level proxy
  - User identities mapped by proxy
- Provides access to user data, VM images

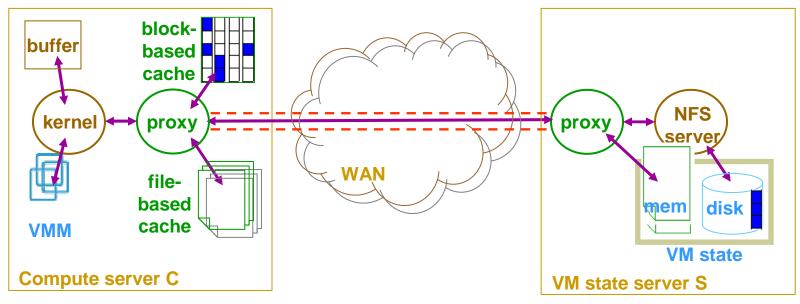


### Many users, apps and environments



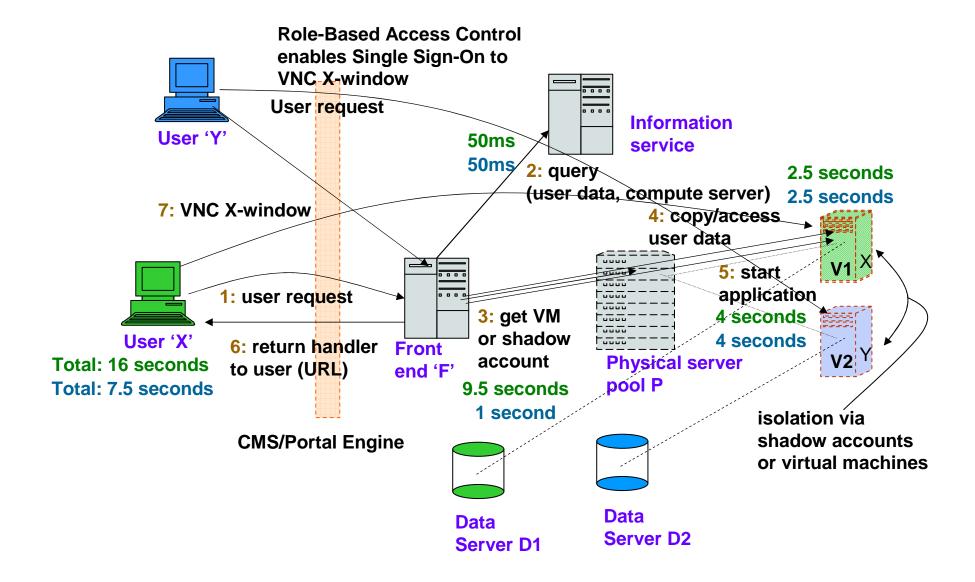
Dynamic, efficient transfer of large VM state is important

## **User-level Extensions**



- Client-side proxy disk caching
- Application-specific meta-data handling
- Encrypted file system channels and cross-domain authentication
- [Zhao, Zhang, Figueiredo, HPDC'04]

### Putting it all together: GUI Application example



## Virtual network services

- VMShop allocates a remote VM
  - Great; now how to access it?
  - Need to isolate traffic from host site
  - For most flexibility, need full TCP/IP connectivity
- ViNe, IPOP being developed at UF ACIS
  - Related work: Virtuoso/VNET (NWU), Violin (Purdue)

### **In-VIGO Virtual Networks - ViNe**

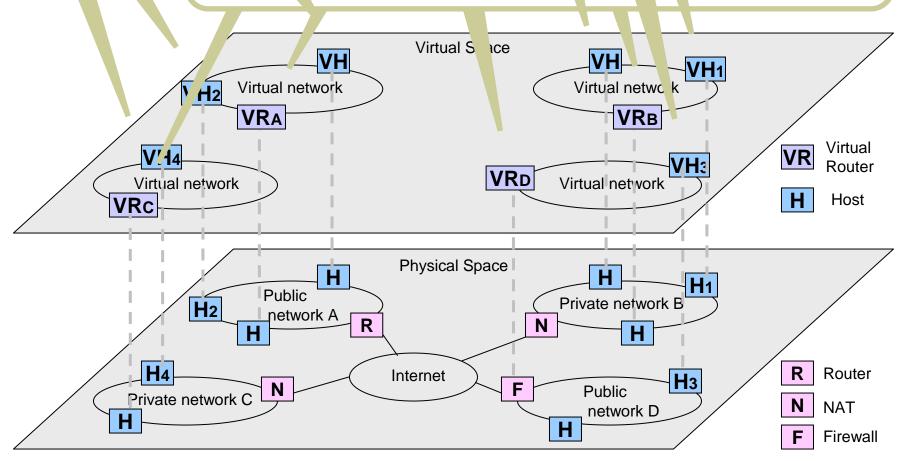
- IP overlay on top of the Internet
- Operation similar to site-to-site VPN
- Designed to address issues that VPN does not solve:
  - High administrative overhead for many sites
  - VPN firewalls need a static public IP address

## **In-VIGO Virtual Networks - ViNe**

- Each participating host is configured with an additional IP address in ViNe space (IP aliasing)
- Packets with destination in ViNe space are directed to VRs for routing in ViNe space.

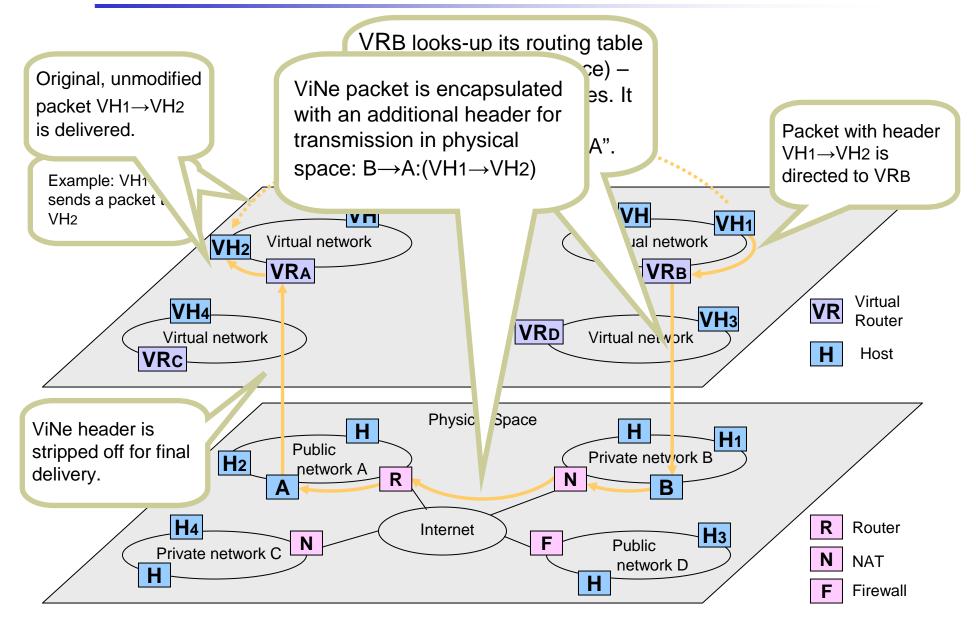
to route (tunnel)

configured with P private space)



Slide provided by M. Tsugawa

## ViNe - Communication in virtual space

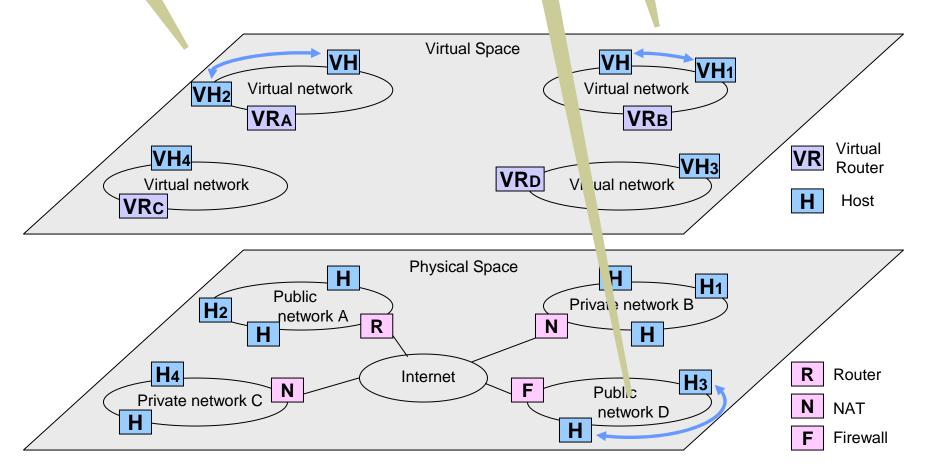


Slide provided by M. Tsugawa



## **ViNe - Local communication**

- Local communication is kept local in both Physical and Virtual space.
- ViNe does not interfere with physical communication.
- Virtual space can be used only when needed.



Slide provided by M. Tsugawa



#### ViNe Firewall/NAT traversal

- VRs connected to the public network proxy (queue) packets to VRs with limited connectivity.
   The latter open connection to the queue VR to retrieve packets.
- VRs with limited connectivity are not used when composing routing tables. Routing tables are made to direct packets to queue VRs.
- The approach supports multi-level NAT.
- The approach also works under DHCP since the changing IP is not considered for routing.

## **ViNe organization**

- Routing tables are created/destroyed as needed (e.g., join/leave of sites, creation of a new ViNe, etc).
- VRs exchange routing information with each other
- Communication of sensitive information (e.g., routing tables, VRs host certificates) is encrypted.
- Administrator of a participating site is involved only during the setup/configuration phase. No intervention is needed when machines join/leave network.

### ViNe Overhead

- When firewall/NAT traversal is not required
  - depends on performance of VRs and available physical network
  - Overhead = 0 5% of available bandwidth
  - up to 150 Mbps for VR on 3 GHz Xeon
- When firewall/NAT traversal is required
  - also depends on the allocation of VRs to proxy/queue traffic
  - 10 50% in initial experiments. Optimizations under investigation.

## **ViNe Security**

#### Site-related

- security policies are not changed by enabling ViNe
  - minimal change may be needed to allow ViNe traffic in private IP space
- ViNe traffic consists of IP packets that are visible in LANs (tunneling is only used across domains)
  - Network policies can be applied to ViNe traffic
  - Firewalls can inspect ViNe traffic
  - Intrusion detection systems and monitoring works unmodified

#### ViNe-related

- ViNe routers do not route packets to/from the Internet
- All communication between VRs are authenticated
- Sensitive VR messages are encrypted
- VRs are not accessible in ViNe space
- ViNe connects hosts without links in physical IP infrastructure
  - But it does so only where we want to have it



# ViNe: On-going work

- Management of Virtual Networks
  - Automated and secure management (definition, deployment, tear-down, merge, split and join/leave of hosts) of virtual networks is under development in the context of ViNe project
  - The idea is to dynamically and securely and reconfigure ViNe routers in response to client (privileged users, local site administrators, grid administrators, grid middleware) requests
  - In collaboration with ANL

# **ViNe: Auditability**

- ViNe does not modify packets generated by participating hosts
- Regular network traffic inspection can be performed in each participating site
- In addition, ViNe Routers can log all routed traffic (performance implications are under investigation)
  - Side-process can combine traffic logs for global network traffic analysis

## **IPOP** virtual network

- Motivations:
  - Enable self-configuring virtual networks focus on making it simple for individual nodes to join and leave
  - Decentralized traversal of NATs and firewalls
- Approach: IP-over-P2P
  - Overhead of adding a new node is constant and independent of size of the network
  - Peer to peer routing
    - Self-organizing routing tables
    - Ring topology with shortcuts
      - N nodes, k edges per node; O(1/k log²(N)) routing hops
      - Adaptive, 1-hop shortcuts based on traffic inspection
  - Mobility: same IP even if VM migrates across domains

A. Ganguly, A. Agrawal, P. O. Boykin, R. Figueiredo - IPDPS 2006, HPDC 2006



# **Applications**

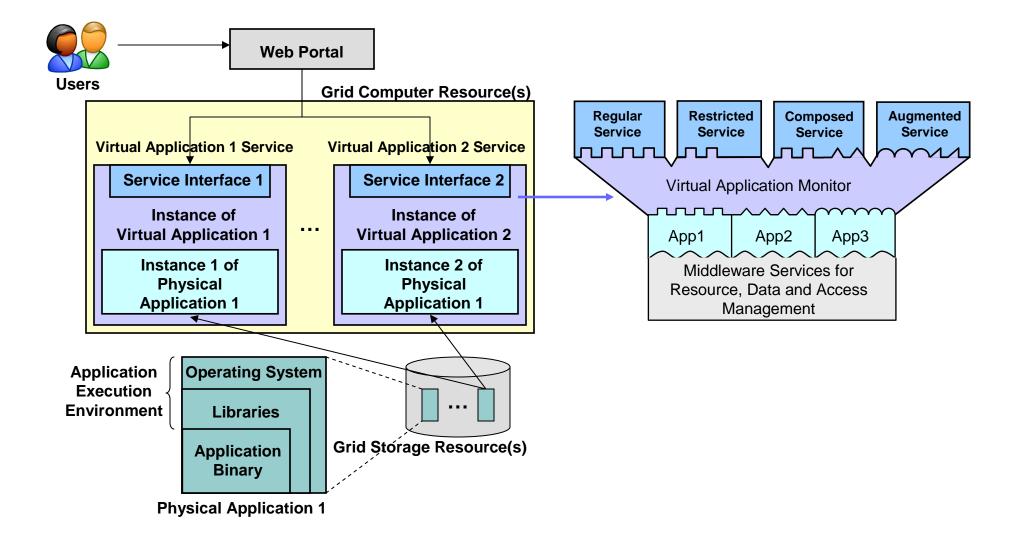
- Distributed computing VM "appliances":
  - Define once, instantiate many
  - Homogeneous software configuration and private network address spaces
  - Facilitates a model where resources are pooled by the various users of a community (e.g. nanoHUB)
  - Homogeneous configuration facilitates deployment of security infrastructures (e.g. X.509-based IPsec host authentication)

# **Usage examples**

- "Grid appliance"
  - Condor node for job submission/execution
  - Automatically obtains a virtual address from virtualized DHCP server and joins a pool
  - Can submit and flock jobs within virtual network
  - Download VMware player and VM image: <a href="http://www.acis.ufl.edu/~ipop/grid\_appliance">http://www.acis.ufl.edu/~ipop/grid\_appliance</a>
- On-going: domain-specific customizations
  - nanoHUB: WebDAV client, Rappture GUI toolkit
  - SCOOP (coastal ocean modeling): clients to access data catalog and archive
  - Archer (computer architecture): support for large, read-only checkpoints and input files

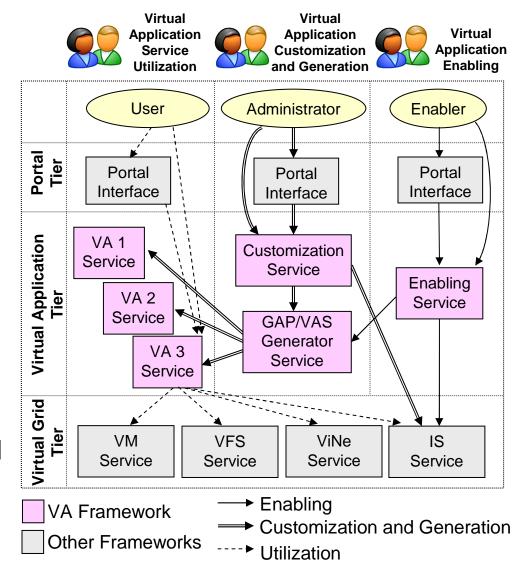


## **Application virtualization**



# **Grid-enabling unmodified applications**

- Enabler provides
  - Command-line syntax
  - Application-related labels
    - Parameter(s), type-set values, entire applications
  - Resource and execution environment metadata
    - Architecture, OS libraries, environment variables
- Grid-services created, deployed and possibly customized using
  - Generic Application Service (GAP)
  - Virtual Application Service (VAS)
- Grid-user interacts with the virtual application through a Web-portal to execute applications on virtualized resources



# **Summary and conclusions**

- Virtualization technology decouples physical resource constraints from user and application requirements
  - Big win, novel rethinking
  - Virtual resources are to grid computing what processes are to operating systems
  - Developers can concentrate on applications, not end resources
- Web-services provide interoperability and a framework for composition and aggregation of applications
  - Includes delivering virtuals and virtualizing applications
  - Wide adoption creates large reusable toolboxes, e.g. for automatic interface generation
  - Users need only know of service interfaces
- In-VIGO middleware effectively integrates virtualization and Web-services technologies to easily enable and deliver applications as Grid-services

# **Acknowledgments**

#### Collaborators

- In-VIGO team at UFL
  - http://www.acis.ufl.edu/invigo
- Rob Carpenter and Mazin Yousif at Intel
- Peter Dinda and Virtuoso team at NWU
  - http://virtuoso.cs.northwestern.edu
- NCN/NanoHub team at Purdue University
  - http://www.nanohub.org
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